

TOWARDS DISTRIBUTED DIAGNOSIS OF REACTIVE SYSTEMS USING STATECHART BASED CONTROLLER DESIGN

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A **fault** is an unexpected change in a system's function and a **failure** is a complete breakdown of a system component or function. **Fault diagnosis** is a process of detecting and isolating faults (fault types): based on event sequences (symptoms). In safety critical applications each possible fault that can result in a symptom is reported as a fault candidate: model based diagnosis is preferred here. In Finite State Machine models of the system to be diagnosed the event set is partitioned as $\Sigma = \Sigma_0 \cup \Sigma_{uo}$: the observable and unobservable event sets. $\Sigma_f \subseteq \Sigma_{uo}$ denotes the set of fault(y) events. Types of faults f_i , ($i=1,p$) can be identified with associated disjoint alphabets Σ_{f_i} , forming a partition, Π_f on Σ_f . A language L is **diagnosable** with respect to Π_f if $(\forall i \in \Pi_f)(\exists n_i \in \mathbb{N})(\forall s \in \Psi(\Sigma_{f_i})) (\forall t \in L/s)(\|t\| \geq n_i \Rightarrow (\forall w \in P_L^{-1}(P(st)))(\Sigma_{f_i} \in w))$; $\Psi(\Sigma_{f_i}) = \{s\sigma_f \in L \mid \sigma_f \in \Sigma_{f_i}\}$; $P: \Sigma^* \rightarrow \Sigma_o^*$; $P_L^{-1}(y) = \{s \in L \mid P(s) = y\}$. This applies only to centralised diagnosis.

However, centralised diagnosis can become impracticable for real-sized systems. Hence, a distributed reference model can be defined as a set of closed languages $\{L_i \subseteq \Sigma_i^* \mid i \in I\}$ where L_i and Σ_i are local components and local event alphabets, respectively. The observable, unobservable and fault event sets for the local component i are, $\Sigma_{io}, \Sigma_{uo}, \Sigma_{if} \subseteq \Sigma_i$, respectively: $i \neq j \Rightarrow \Sigma_{io} \cap \Sigma_{jo} = \Phi$; $i \neq j \Rightarrow \Sigma_{if} \cap \Sigma_{jf} = \Phi$. The distributed diagnosis problem is then posed as, Local computation: $(\forall i \in I) M_i := P_{io}^{-1}(u_i) \cap L_i$: given observation u_i , ($i \in I$) find all strings that can exhibit u_i ; and Global consistency: $E := Sup\Delta(\{M_i \mid i \in I\})$: given the local estimates, use the notion of supremal global support to capture "agreement" among them. Checking for global consistency can be very time consuming. A set of local languages $L = \{L_i \subseteq \Sigma_i^* \mid i \in I\}$ is globally consistent if $\forall i \in I, L_i = P_{i,i}(\prod_{j \in I} L_j)$. Here, $P_{j,i}: \Sigma_j^* \rightarrow (\Sigma_j \cap \Sigma_i)^*$ and \prod represents the synchronous product of languages. L is locally consistent if $\forall i, j \in I, P_{i,j}(L_i) = P_{j,i}(L_j)$. Global consistency (GC) implies local consistency(LC), but not vice-versa. Our distributed diagnosis strategy uses conditions that make this latter equality possible: let graph Gr (Ver, Edg) be constructed with local event sets as vertices ($Ver = I$) and edges connecting event sets with non-disjoint alphabets. Then, Gr is a tree $\Rightarrow \{LC \Rightarrow GC\}$. We were also inspired by a statechart based approach using structures called D-Holons to represent superstates with associated diagnosers built as Reachability Transition Systems (RTS) of observation-adjacent states. However, their approach is very restrictive in that they only consider **failures** and also single-failure scenarios. We manage to overcome the restrictions by adopting Drusinski-Harel decomposition on statechart models of the system to be controlled using an approach we have developed for modular verification of Logic/supervisory control: this decomposition satisfies the "Tree" condition required for the graph Gr, to make $LC \Rightarrow GC$.